



On Colliding First Messages in Slotted ALOHA

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Motivation

Several techniques in wireless and wired networking require some method for distributed node selection.



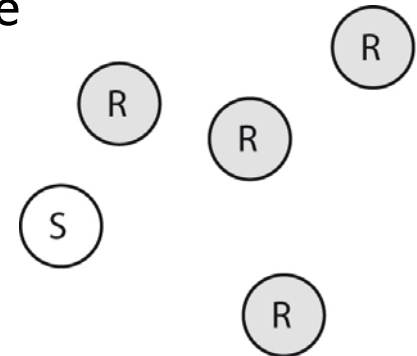
Examples:

- Cooperative relaying techniques in wireless networks to choose a “relay node”
- Data processing techniques in sensor networks to choose a “data gathering node”

A possible way to perform node selection

Step 1: Determine a set of candidate nodes

- A node broadcasts a query message to all neighboring nodes.
- This message indicates a certain criterion (or several criteria) that qualifies to serve as a selected node
- Each receiving node that fulfills the criterion becomes part of a candidate set.

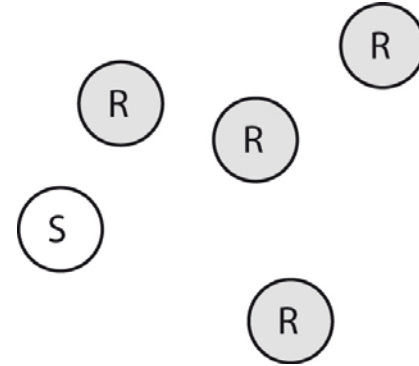


Step 2: Determine the selected node

- All nodes of the candidate set compete for random access on the shared medium, to send back a reply.
- The node that successfully accesses the medium *first* wins the selection process and acts as selected node.

The importance of “first messages”

- The reply message of the node answering *first* is more *important* than subsequent reply messages.
- This reply message should *not collide* with other messages, hence only one node should access the channel.



This discussion leads to the following **MAC design issues**:

- What is the probability that there is a first message that does not collide?
- How can we maximize this probability?
- What is the tradeoff between this probability and the delay of the selection process?

Modeling assumptions and definitions



- n devices
- s slots with slotted ALOHA

Design parameter: Each device transmits with probability p_i in slot i

- If two or more nodes transmit in the same slot, a *message collision* occurs. A message not suffering from a collision is called a *non-colliding message*.
- A slot is *empty* if no node transmits during this slot. The *first non-empty slot* is the slot i in which at least one message is sent while previous slots $1, \dots, i-1$ were empty. A message sent in the first non-empty slot is called a *first message*.

Non-colliding first message probability

What is the probability that there occurs a non-colliding first message within s slots?

$$\Phi(n, s, p_1, \dots, p_s) =$$

$$= n p_1 (1 - p_1)^{n-1} + (1 - p_1)^n \cdot n p_2 (1 - p_2)^{n-1} + \dots$$

Exactly 1 message in slot $i = 1$. *or* No message in slot $i = 1$. *and* Exactly 1 message in slot $i = 2$.

$$= n \sum_{i=1}^s \left(\prod_{w=0}^{i-1} (1 - p_w)^n \right) \cdot p_i (1 - p_i)^{n-1} \quad \text{with } p_0 := 0.$$

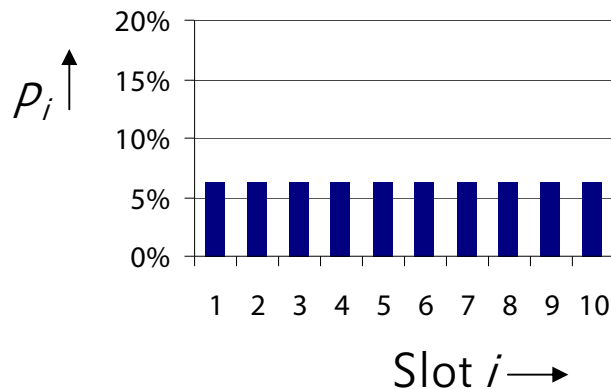
Optimizing the non-colliding first message probability

How to set the transmission probabilities p_i to maximize the probability Φ of obtaining a non-colliding first message within s slots?

Example: $n = 5$ nodes on a channel with $s = 10$ slots

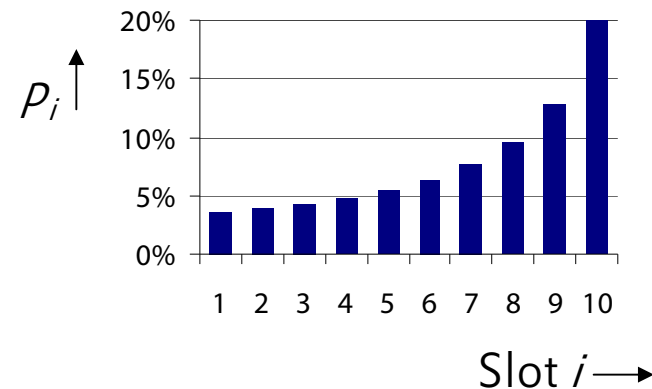
(a) Each slot i has same p_i :

$\Phi_{\max} = 84.05\%$ obtained with



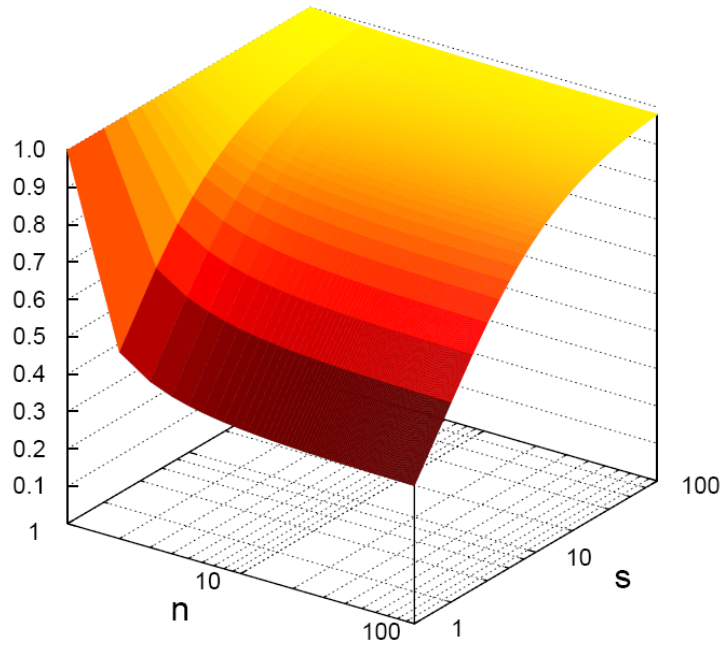
(b) Each slot i may have different p_i :

$\Phi_{\max} = 86.68\%$ obtained with

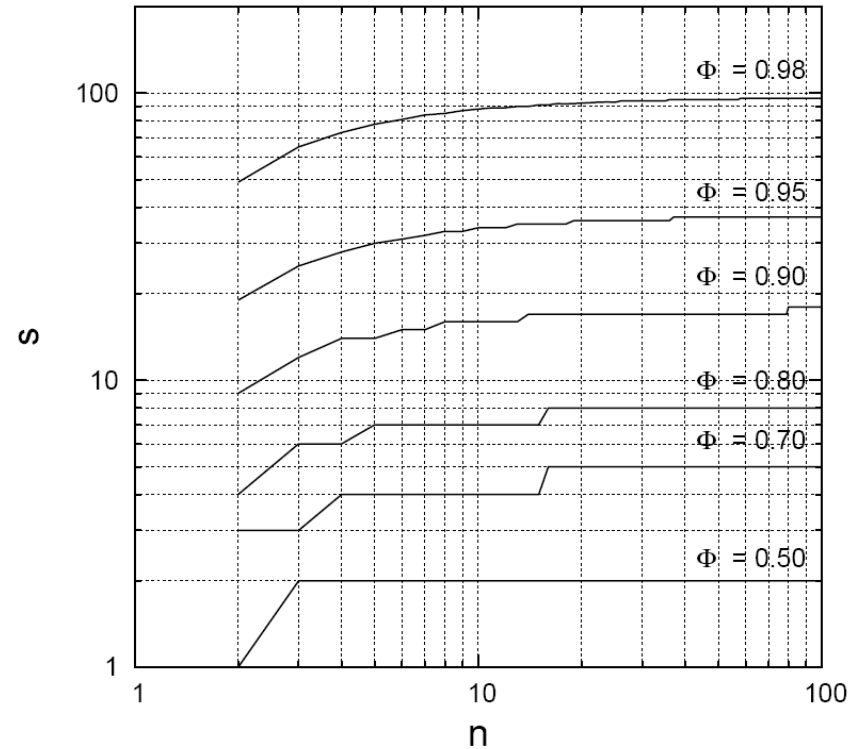


Slow start strategy

Maximum possible non-colliding first message probability



(a) $\max \Phi(n, s, p_1, \dots, p_s)$



(b) Contour lines of $\max \Phi(n, s, p_1, \dots, p_s)$

Optimizing the non-colliding first message probability

Probability Φ gets maximum if we set the transmission probabilities to

$$p_{s-k} = \frac{1}{n} \cdot \frac{(n-1)^k - n \alpha_k \beta_k}{(n-1)^k - \alpha_k \beta_k} \quad \text{with index } k \in \{0, \dots, s-1\}$$

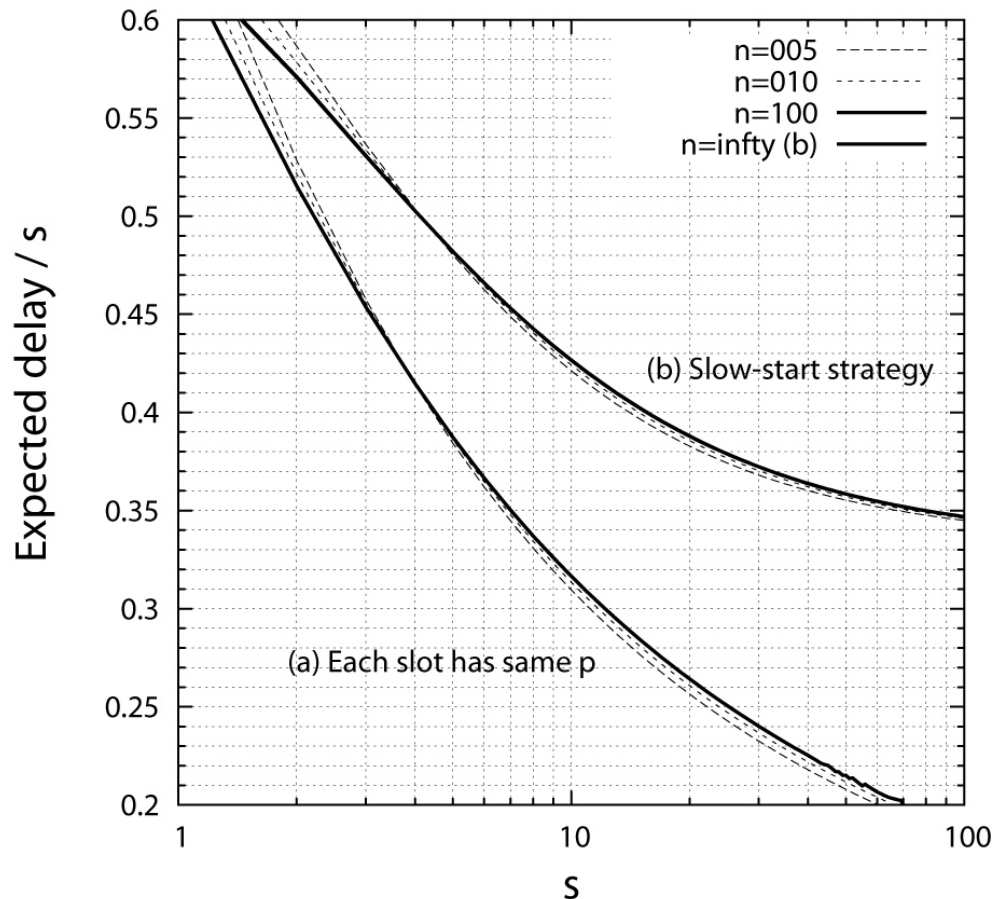
$$\alpha_k := \left(\frac{(n-1)^k}{n} \right)^n \quad ; \quad \beta_k := \begin{cases} 0 & \text{for } k = 0 \\ ((n-1)^{k-1} - \alpha_{k-1} \beta_{k-1})^{1-n} & \text{else} \end{cases}$$

Important observation:

- A node is not forced to transmit within s slots ($\sum p_i \neq 1$)
- If we force each node to transmit within s slots, a worse probability is obtained.

Delay of the first message

What is the expected delay of the first message? What is the maximum delay that can be guaranteed in 90% of all cases?



Related work

- T. Watteyne, I. Augé-Blum, M. Dohler, D. Barthel: “Reducing collision probability in wireless sensor network backoff-based election mechanisms.” In *Proc. IEEE GLOBECOM*, (Washington, DC), Nov. 2007.
- Y. Tay, K. Jamieson, H. Balakrishnan: “Collision-minimizing CSMA and its applications to wireless sensor networks.” *IEEE J. Select. Areas Commun.*, Aug. 2004.
- J. A. Stine, G. de Veciana, K. H. Grace, R. D. Durst: “Orchestrating spatial reuse in wireless ad hoc networks using synchronous collision resolution.” *J. Interconnection Networks*, Sept. 2002.

Conclusions and outlook

Analytical analysis of “non-colliding first messages” on a link with n nodes performing random access using ALOHA with s slots.

- The probability Φ that there occurs a non-colliding first message within the given slots is maximized by a **slow start strategy** of the nodes.
- We can calculate this optimal probability Φ and the sending probabilities p_i leading to this Φ , if we know n and s .
- The slow start strategy comes at the price of an **increased delay** of the first message; this delay is almost independent of n .

Outlook

- The number of nodes n is not always known. What is the sensitivity of Φ with respect to n ? How can we estimate it?

